

Fuzzy Presheaves are Quasitoposes

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joint work with Roy Overbeek and Jörg Endrullis

12 May 2023, **VU TCS Seminar**

Vrije Universiteit Amsterdam



Fuzzy sets

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Definition

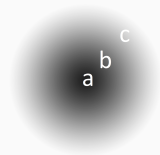
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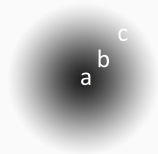
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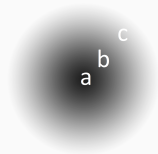
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More general definition:

poset (\mathcal{L}, \leq) instead of $[0, 1]$.



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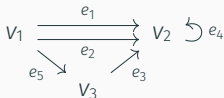
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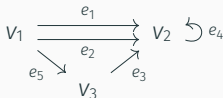
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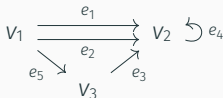
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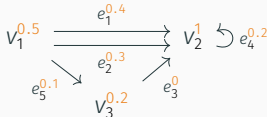
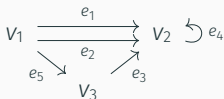
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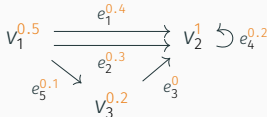
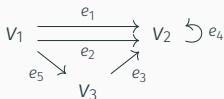
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For some authors: $v_1^x \xrightarrow{e^z} v_2^y \Rightarrow x, y \geq z$.

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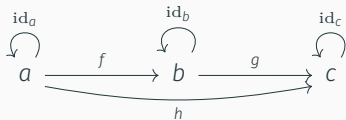
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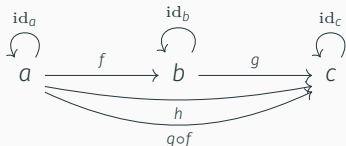


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- composition of arrows $a \xrightarrow{g \circ f} b, \quad \dots$



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- VectSpace: category of *vector spaces* and *linear transformations*

Graphs as Presheaves

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 $\implies A$ is a **graph**.

More examples

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$$\blacksquare I^{\text{op}} = E \xleftarrow{f} R \xrightarrow{g} V. \text{ Given a presheaf } G, \text{ read } r \in A(R) \text{ as meaning}$$

“vertex $A(g)(r) \in A(V)$ is incident to edge $A(f)(r) \in A(E)$ ”.

\implies multiple incidences hypergraphs

More examples (continued)

Let us come back to 2-element edges.

- **Undirected graphs:** $e = (v_1, v_2)$, $\text{sym}(e) = (v_2, v_1) \iff e = \{v_1, v_2\}$

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- **Undirected reflexive:** Both sym and refl . Extra equation: $\text{sym} \cdot \text{refl} = \text{refl}$.

Graphs homomorphisms as natural transformations

What about graph homomorphism?

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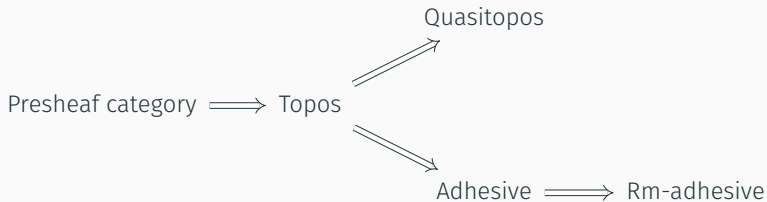
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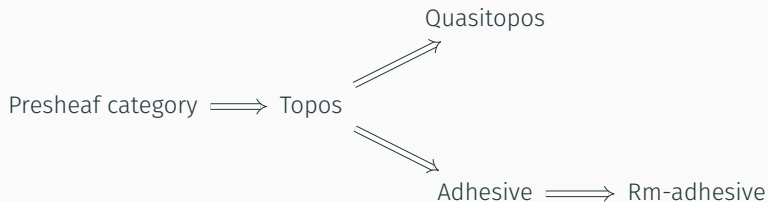
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Thus, $\text{Presheaf}(I) \cong \text{Set}$.
- For $I^{\text{op}} = E \begin{matrix} \xrightarrow{s} \\ \xrightarrow{t} \end{matrix} V$, then $f : A \Rightarrow B$ is a *graph homomorphism*:
 - ▶ two function $f_E : A(E) \rightarrow B(E)$ and $f_V : A(V) \rightarrow B(V)$
 - ▶ such that $B(s) \circ f_E = f_V \circ A(s)$ and $B(t) \circ f_E = f_V \circ A(t)$

Thus $\text{Presheaf}(I) \cong \text{Graph}$.

Presheaves properties



Presheaves properties



Recall (last TCS seminar) [Overbeek et al, 2023]:

- in quasitopos: PBPO^+ subsumes DPO, SqPO, AGREE, PBPO
- in rm-adhesive quasitopos: termination method using weighted subgraph counting.

Fuzzy Presheaves

Why care about fuzzy structure?

Membership values can be seen as labels $v_1 \xrightarrow{e^x} v_2$.

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When (\mathcal{L}, \leq) is a complete lattice (i.e., join \vee , meet \wedge , bot \perp , top \top), PBPO⁺ offers an easy treatment of *graph relabelling*, i.e., changing a label x (of an edge or vertex) to another label y

Presentations of monads by algebraic theories

Definition (Informally)

A **fuzzy presheaf** is a presheaf $A : I^{\text{op}} \rightarrow \text{Set}$ where every set $A(i)$ has a membership function $\alpha_i : A(i) \rightarrow \mathcal{L}(i)$ to some poset $(\mathcal{L}(i), \leq)$.

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(A **fuzzy presheaf morphism** is a natural transformation $f : (A, \alpha) \rightarrow (B, \beta)$ such that $\alpha_i \leq \beta_i f_i$ for all $i \in I^{\text{op}}$.)

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Presentations of monads by algebraic theories

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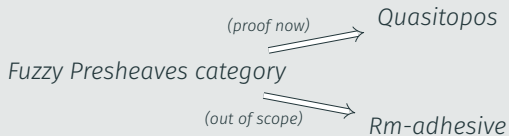
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Quasitopos

Theorem

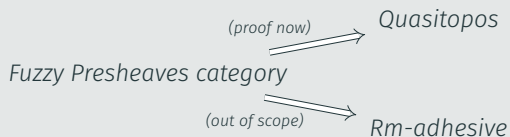
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To be a quasitopos, three things needed:

- finite limits, finite colimits
- regular-subobject classifier
- exponent objects (roughly)

Finite (co)limits

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Lemma

Fuzzy presheaf $1(i) := \{\cdot^1\}$ for all $i \in I$ is a terminal object.

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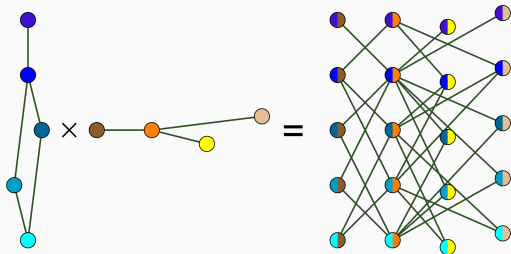
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Last limits

There are other limits to verify: equalizers or pullbacks, but it is similar.

Finite colimits are dual:

- initial object
- coproduct (i.e., disjoint union)

Classifier Ω

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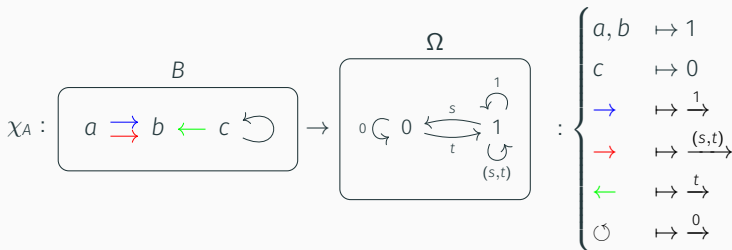
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Regular-subobject classifier

- Let $\Omega := \{0^1, 1^1\}$. Fuzzy subsets $A_1 = \{a^{0.1}\}$, $A_2 = \{a^{0.4}\} \subseteq B = \{a^{0.4}, b^{0.6}\}$.

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Lemma

*The subobject classifier Ω for presheaves, with every element given full membership 1 is a **regular**-subobject classifier for fuzzy presheaves.*

Exponential

Definition

For objects A, B of a category, an **exponential object** is an object B^A s.t.

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Future work: Define “fuzzy quasitopos”? Prove they are quasitoposes?